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# Crowdsensing Data Trading for Unknown Market: Privacy, Stability, and Conflicts

He Sun<sup>1</sup>, Mingjun Xiao<sup>1</sup>, Yin Xu<sup>1</sup>, Guoju Gao<sup>2</sup>, Shu Zhang<sup>1</sup>



IEEE TRANSACTIONS ON  
MOBILE COMPUTING

## Key Question

How to design a Crowdsensing Data Trading Framework considering **privacy** and **stability** for **unknown market** in **centralized**<sup>[1]</sup> and **decentralized**<sup>[2]</sup> settings?

## Introduction

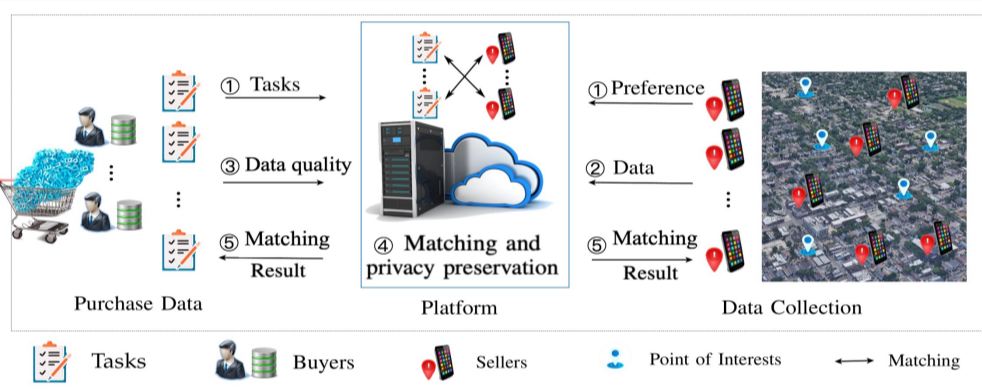
### ➤ Crowdsensing Data Trading (CDT)

A new data trading paradigm where the Mobile CrowdSensing (MCS) technique is adopted to provide data sources, e.g., Thingful, ThingSpeak.

### ➤ Concept of Matching Markets

- ✓ Both sides of the markets can't just choose what you want even if you can afford it.
- ✓ One of them also have to be chosen.
- ✓ They choose each other according to the preferences of each other.

### ➤ Components of CDT systems



### PS-CDT platform

**Platform:** As a broker, it provides credible data trading services for sellers and buyers.

**Buyers:** Propose and publish their data requirements to the platform to collect data.

**Sellers:** A crowd of mobile users to provide data collection service to buyers.

### ➤ Existing Problems

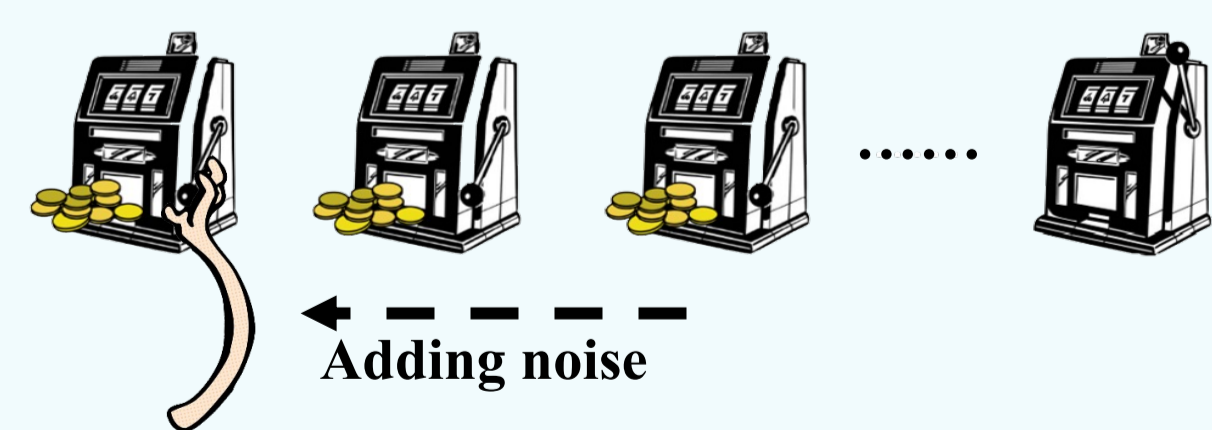
- A few existing CDTs consider the stability of the Data Trading Market.
- The Data Trading Market is unknown in practice, i.e., the preference sequences over sellers are unknown by buyers.
- The private information of sellers needs to be preserved.
- Decentralized CDT has more practical significance.
- Two matching requests for the same seller would create a competitive matching conflict.

### ➤ Contributions

- ✓ To the best of our knowledge, this is the first CDT work that takes the unknown market, privacy preservation, and the stability of the data trading into consideration simultaneously in centralized and decentralized settings.
- ✓ We define a novel metric, i.e.,  $\delta$ -stability to measure the stability of the markets.
- ✓ We propose the DPS-CB and CDPS-CB mechanisms to solve the privacy, stability, and conflicts-avoiding problems.

## System, Modeling, and Problem

### ➤ $\epsilon$ -Differentially private bandit model



Platform	Game players
Sellers	Arms
Select a seller	Pull an arm
Data quality	Reward
Protected Data quality	Perturbed reward

$$\mathbb{P}\{\Phi_i(q_i^{1:l-1}) \in \mathcal{X}\} \leq e^\epsilon \cdot \mathbb{P}\{\Phi_i(q_i^{1:l-1}') \in \mathcal{X}\} \quad (1)$$

✓ where  $\epsilon > 0$  is a small constant that the policy provides, indicating the privacy-preserving level.

✓ A bandit policy  $\Phi_i$  of play  $i$  is a sequence of arm-pulling decisions.

✓  $q_i^{1:l} = \{q_i^1, \dots, q_i^l\}$ ,  $q_i^{1:l-1}$  is its adjacent sequence.

### ➤ $\delta$ -Stable Matching Model

#### Definition of preference

*Unknown preference sequences of the buyer*

*Preference sequence of the seller*

✓ Denoted by  $\pi_k^l = \{\dots, \pi_k^l, \dots\}$ ,  $\pi_k^l = \{\dots, j, j', \dots\}$ .

✓ Denoted by  $\pi_j = \{\dots, i, i', \dots\}$ .

✓  $\pi_k^l(j)$  denotes the rank of seller  $j$  in  $\pi_k^l$ .

✓  $\pi_j(i)$  denotes the rank of task  $i$  in  $\pi_j$ .

✓  $v_i = \{\dots, v_i^l(j), \dots\}$  denotes the value. (**Unknown**)

Adding noise  $\rightarrow$  Matching is **not truly** stable

**Definition of  $\delta$ -stable:** We say a market outcome  $M^l$  is  $\delta$ -stable with a probability less equal than  $1 - \delta$  that a preference sequence is invalid, i.e., there exists two matching pairs  $\langle i, j \rangle$  and  $\langle i, j^* \rangle$ ,  $\forall i \in T, \forall j, j^* \in S$ , satisfies  $\pi_k^l(j) < \pi_k^l(j^*)$ ,  $\hat{\pi}_k^l(j) < \hat{\pi}_k^l(j^*)$  and  $\hat{v}_k^l(j) - \hat{v}_k^l(j^*) > \xi'_0$ , denoted by  $\tilde{M}^k$ .  $\xi'_0$  is a perturbed care bound and  $\delta$  is a constant less than but close to 1.  $\pi_k^l(j) < \pi_k^l(j^*)$ : task  $i$  prefers seller  $j$  to  $j^*$  in  $l^{\text{th}}$  round.

### ➤ Problem formulation

Our **goal** is to **make the optimal matching** in each round according to the built perturbed preference sequences, i.e., to maximize the expected accumulative reward for each task, **assuring the  $\epsilon$ -differential privacy** and  **$\delta$ -stable of market outcomes** in each rounds.

$$\text{Maximize: } \sum_i q_i^l(m^l(i))$$

Subject to: Eq. (1) holds

$M^l$  is  $\delta$ -stable

## DPS-CB and CDPS-CB mechanisms

#### Algorithm 1: DPS-CB mechanism

**Input:** the total rounds  $N$ , the preference sequences set  $\{\pi_j | \forall j \in S\}$  of sellers.  
**Output:**  $\{M^l | l = 1, 2, \dots\}$

- for  $l = 1, \dots, N$  do
- if  $l \leq T$  then
- $m^l(i) \leftarrow j$ ,  $\forall j \in S$ ;
- Get  $q_i^l(j)$  as the corresponding reward according to Eqs. (6-8) while using  $\epsilon$  as the privacy budget under the hybrid differentially private mechanism;
- else if  $l = T + 1$  then
- Compute the DP-UCB indexes  $I_i^l(j)$ ,  $\forall i \in T$ ,  $\forall j \in S$  according to Eq. (9);
- Sort the sellers by the DP-UCB index to build the initial perturbed preference sequence  $\hat{\pi}_i^l$  of each task over sellers;
- Compute stable matching to get the market outcome  $M^l$  according to  $\{\pi_j | \forall j \in S\}$  and  $\{\hat{\pi}_i^l | \forall i \in T\}$  using the Gale and Shapley algorithm;
- else
- Update  $I_i^l(j)$ ,  $\forall i \in T$ ,  $\forall j \in S$  and  $\{\hat{\pi}_i^l | \forall i \in T\}$  according to Eqs. (6-9).
- Compute stable matching to get the market outcome  $M^l$  in the way of Step 8.
- end
- end

#### Algorithm 2: CDPS-CB mechanism

**Input:** the preference sequences set  $\{\pi_j | \forall j \in S\}$  of sellers, the Bernoulli mean  $p$   
**Output:**  $\{M^l | l = 1, 2, \dots\}$

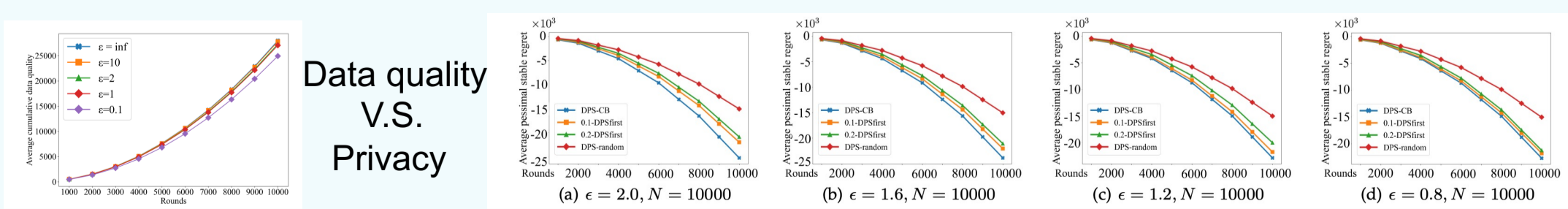
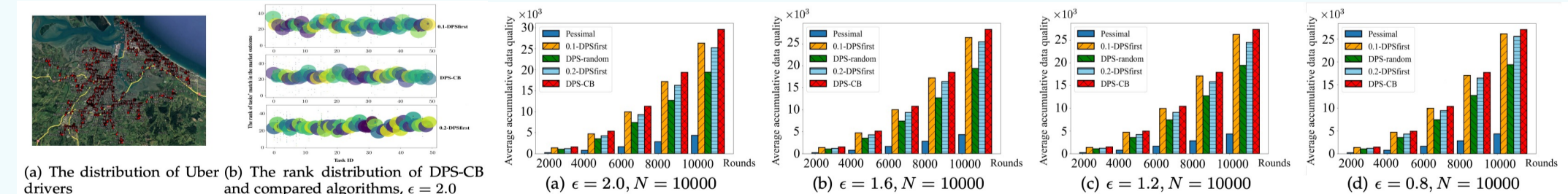
- Initialization:**
- $I_i^l(j) = +\infty$ ,  $\forall i \in T$ ,  $j \in S$ ;
- Find a matching that one-to-one maps from  $i$  to  $j$  randomly,  $\forall i \in T$ ,  $j \in S$ ;
- for  $l = 0, \dots, N$  do
- for  $i = 1, \dots, T$  do
- Sample a random value  $B^l(i)$  from  $Ber(p)$ ;
- if  $B^l(i) = 0$  then
- Update the set of feasible sellers according to Eqs (10);
- Task  $i$  selects the seller with the maximum DP-UCB index to match:
- $m^l(i) = \max\{I_i^l(j) | \forall j \in F^l(i)\}$ ;
- end
- else
- Task  $i$  matches the same seller as the last round:
- $m^l(i) = m^{l-1}(i)$ ;
- end
- if  $i$  wins the conflicts then
- $M^l \leftarrow \langle i, m^l(i) \rangle$ ;
- Update  $I_i^l(m^l(i))$  and  $\{\hat{\pi}_i^l | \forall i \in T\}$ ;
- end
- end
- end

### ➤ Theoretical Analysis

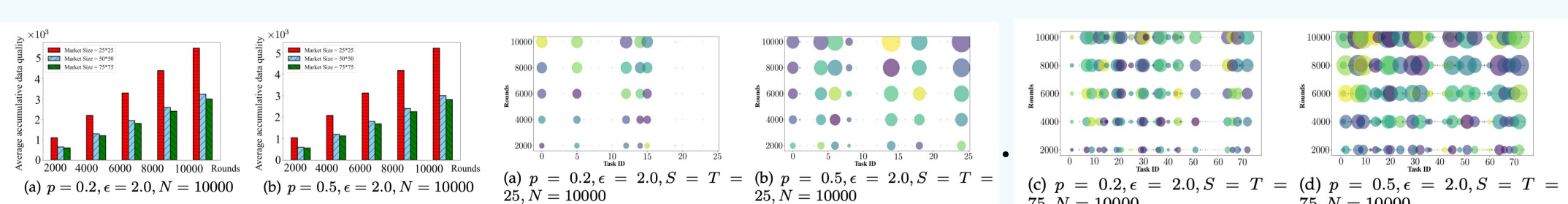
- ✓ The two mechanisms satisfies  $\epsilon$ -differential privacy.
- ✓ The market outcome computed by twos mechanisms are  $\delta$ -stable.
- ✓ The two mechanisms can achieve **sublinear pessimal stable regret**

## Performance Analysis

### ➤ The performance of DPS-CB mechanism.



### ➤ The performance of CDPS-CB mechanism.



[1] He Sun, Mingjun Xiao, Yin Xu, Guoju Gao, Shu Zhang "Privacy-preserving Stable Crowdsensing Data Trading for Unknown Market", IEEE INFOCOM'23, May. 2023

[2] He Sun, Mingjun Xiao, Yin Xu, Guoju Gao, Shu Zhang. "Crowdsensing Data Trading for Unknown Market: Privacy, Stability, and Conflicts", IEEE Transactions on Mobile Computing, 2024.